



# Automated Program Analysis: Revisiting Precondition Inference through Constraint Acquisition

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# Speaker



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# On the Way to Secure Code



## Improve Confidence in Software

↳ Testing

↳ Formal Verification

– E.g., Precondition / postcondition

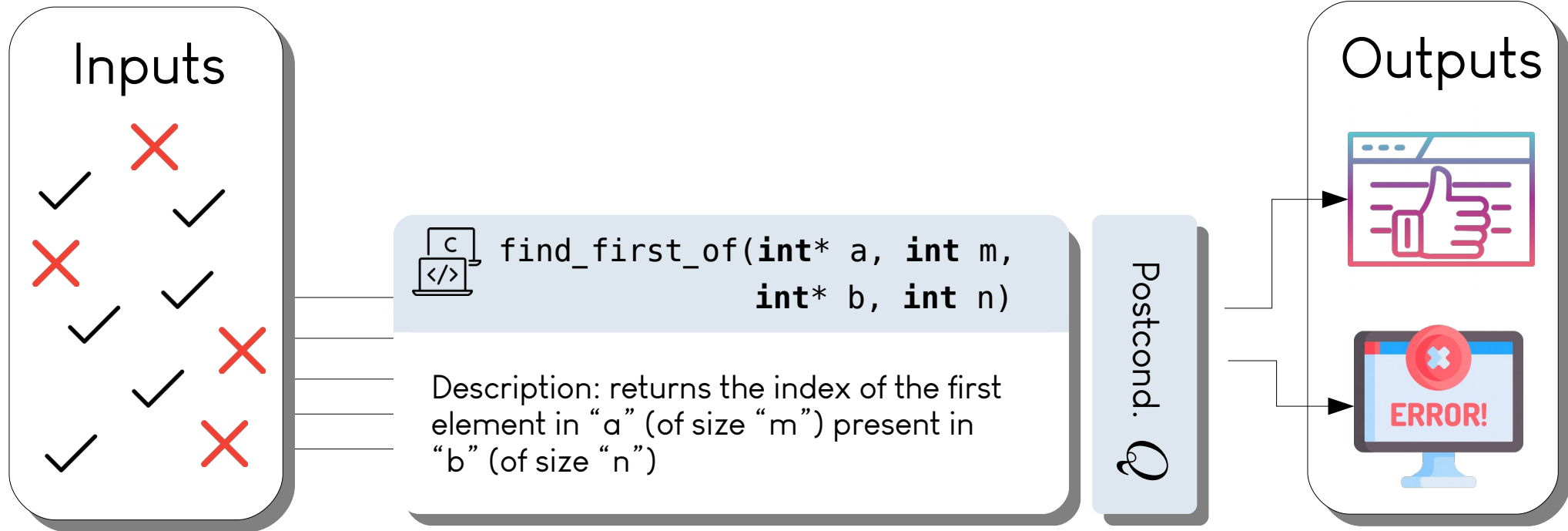


Enable to scale to big code

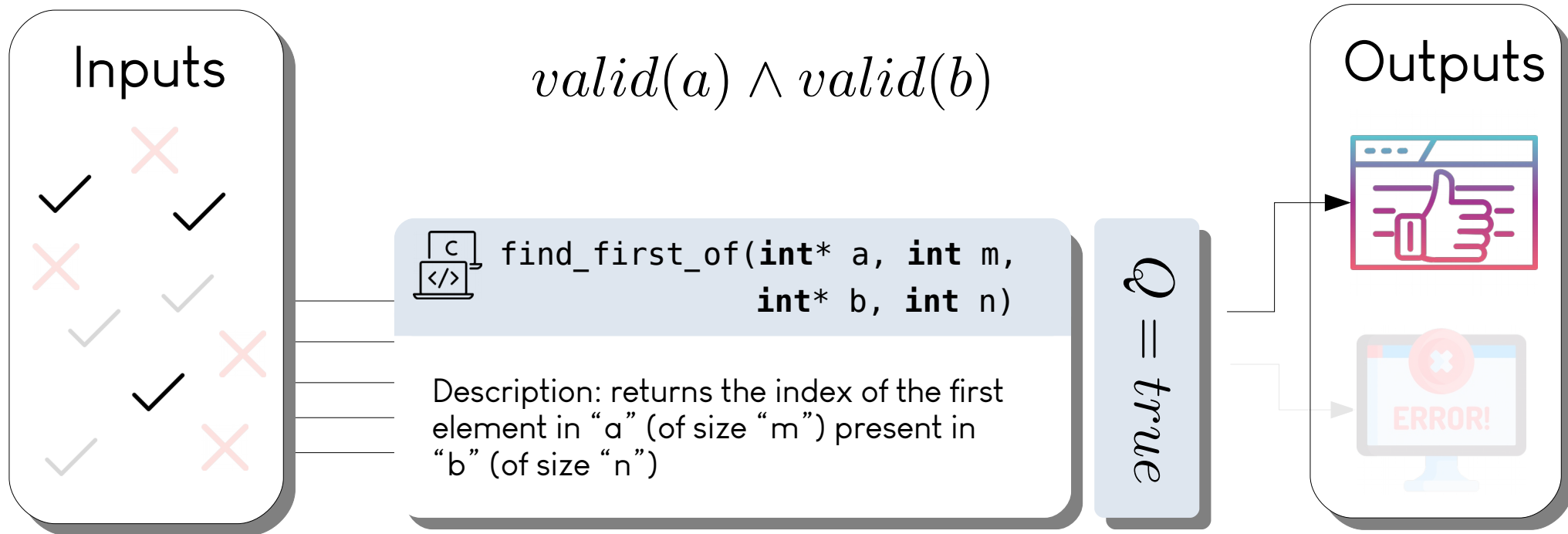


Almost never given in practice

# Dream: Infer Preconditions

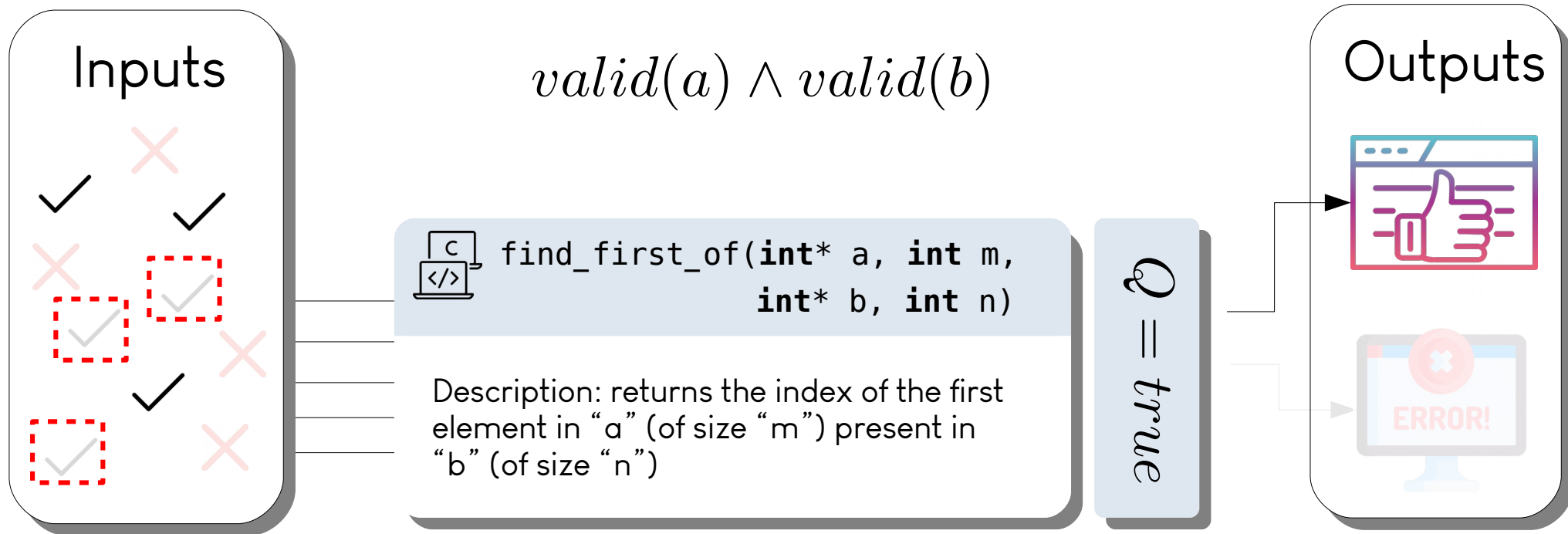


# Dream: Infer Preconditions



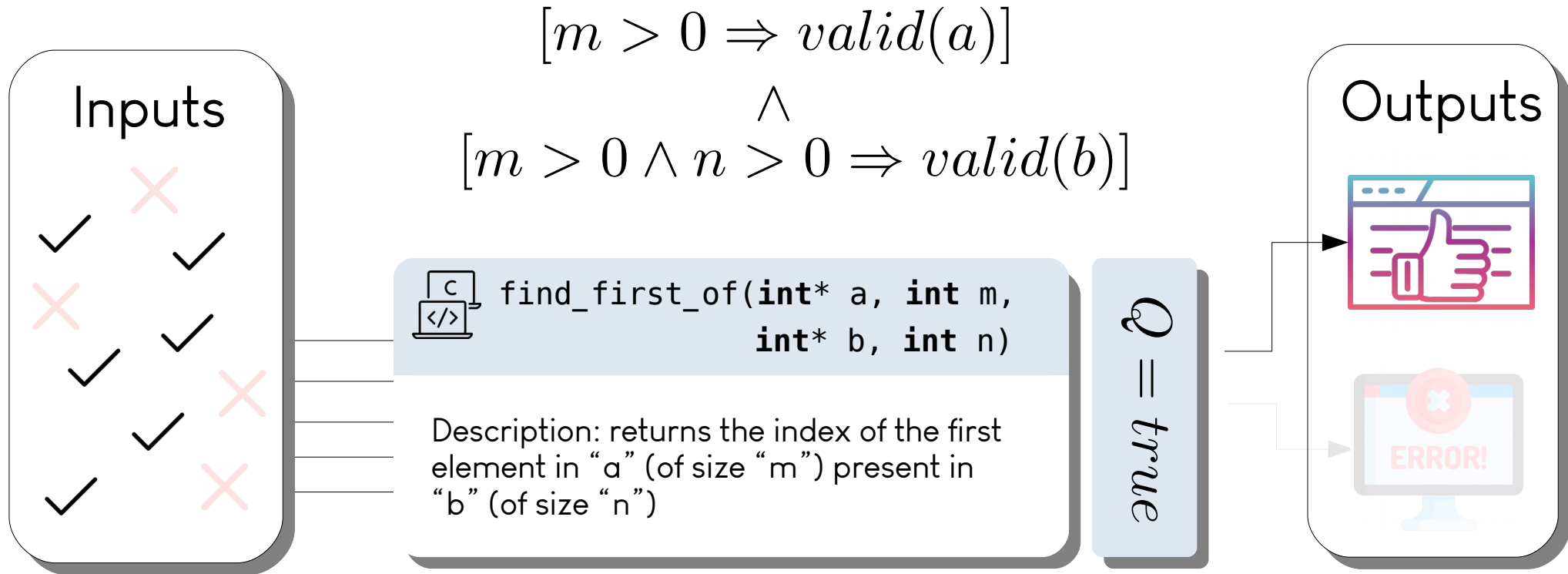
Undecidable problem: Rice theorem (1953)

# Dream: Infer Preconditions



Undecidable problem: Rice theorem (1953)

# Dream: Infer The Weakest Precond.



Undecidable problem: Rice theorem (1953)



# State-of-the-art

Execution Based (Daikon, PIE, Gehr et al.):



Does not need the source code



No clear guarantees

Code Based:



Need the source code

– scalability issues • code not available



Clear guarantees

## Data-Driven Precondition Inference with Learned Features

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## Counterexample-Guided Precondition Inference\*



Mohamed Nassim Seghir and Daniel Kroening  
Computer Science Department, University of Oxford



# Goal





## Execution Based (Daikon, PIE, Gehr et al.):

-  Does not need the source code
-  Clear guarantees

Constraint Acquisition  
Based Precond.  
Inference

## Code Based:

-  Need the source code
  - scalability issues • code not available
-  Clear guarantees

### Data-Driven Precondition Inference with Learned Features

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# Constraint Acquisition



## Constraint Programming

↳ Hard to design models

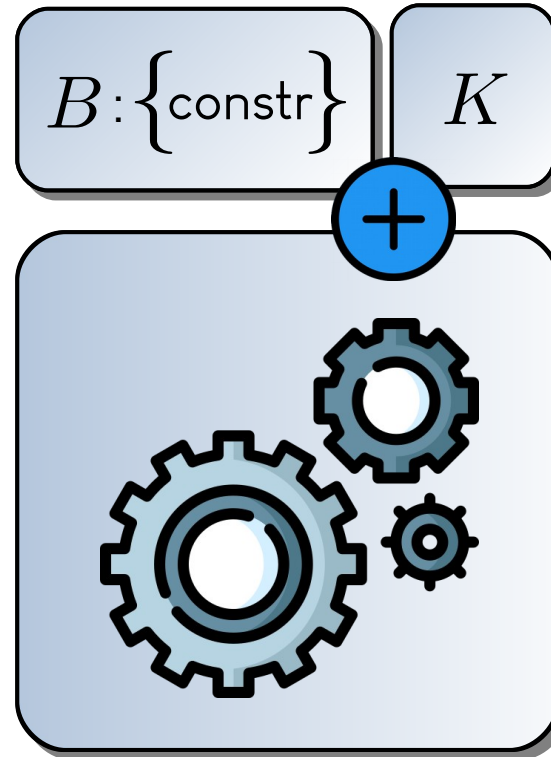


## Constraint Acquisition

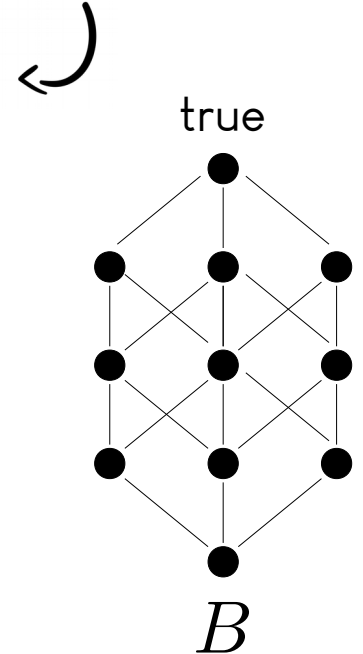
↳ Version Space Learning (Mitchell, 82)

↳ Bessiere, C., Koriche, F., Lazaar, N., & O'Sullivan, B. (2017).  
Constraint Acquisition. *Artificial Intelligence*, 244, 315-342.

# Active Constraint Acquisition



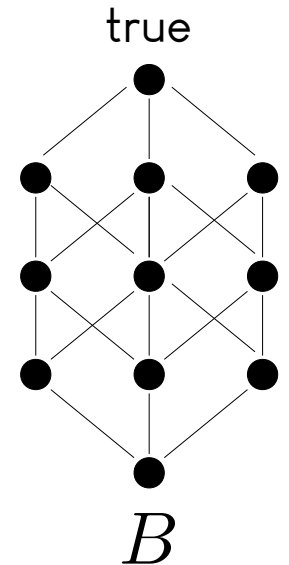
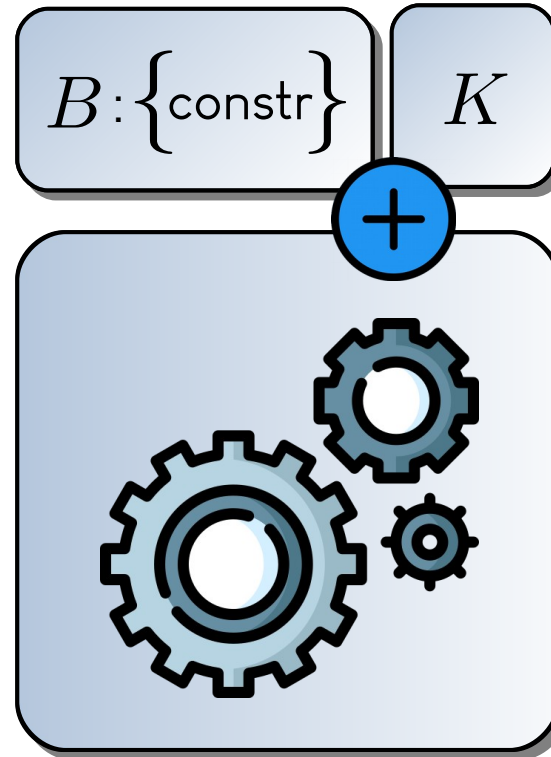
Background knowledge:  
rules to speed up learning



# Active Constraint Acquisition



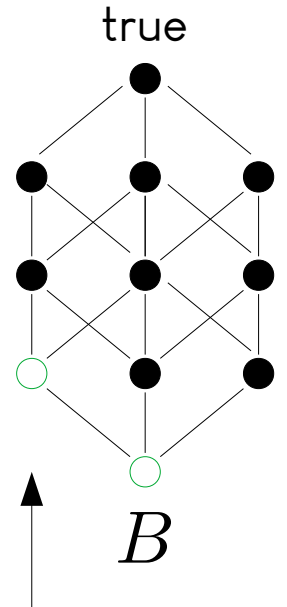
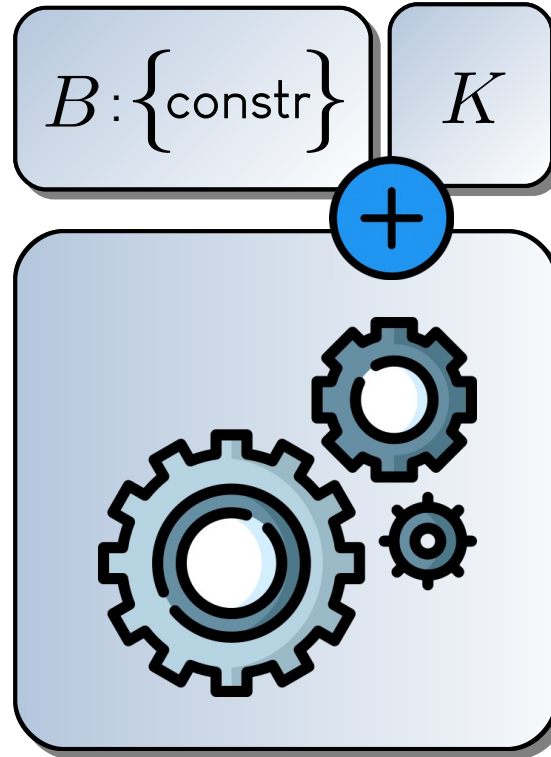
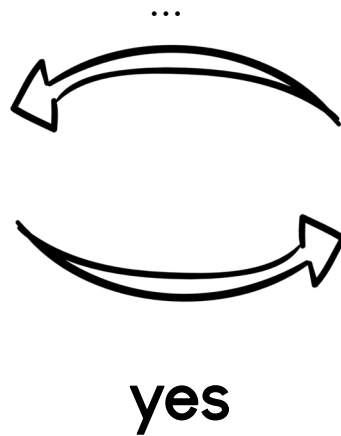
Query  
Elise: 8h - 12h  
Paul: 10h - 11h



# Active Constraint Acquisition



Query  
Elise: 8h - 12h  
Paul: 10h - 11h

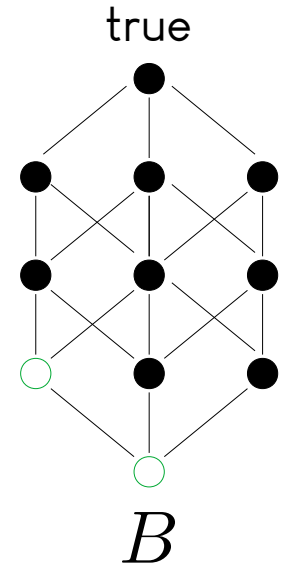
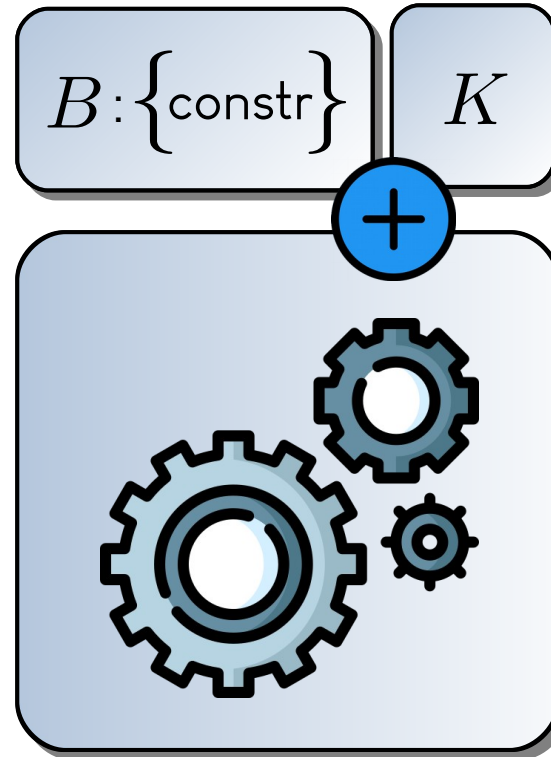


yes: Bottom-up

# Active Constraint Acquisition



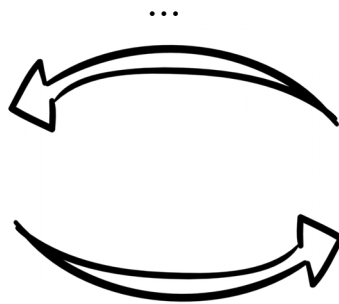
Query  
Elise: 8h - 12h  
Paul: 10h - 23h



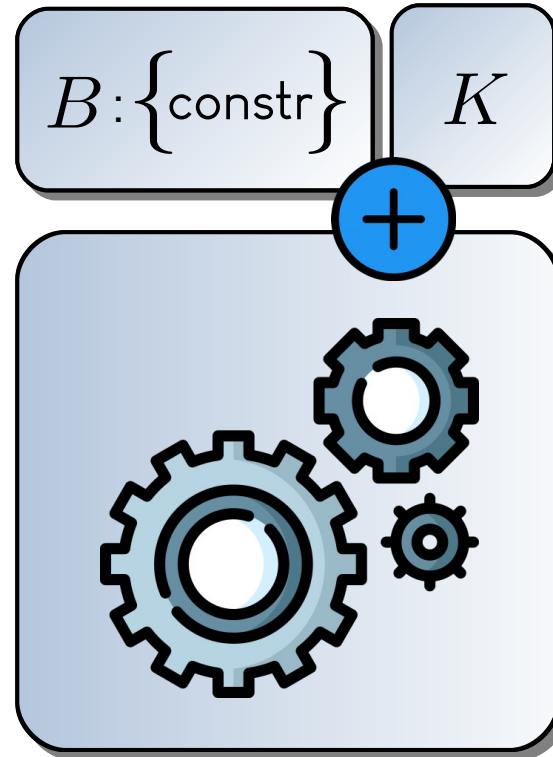
# Active Constraint Acquisition



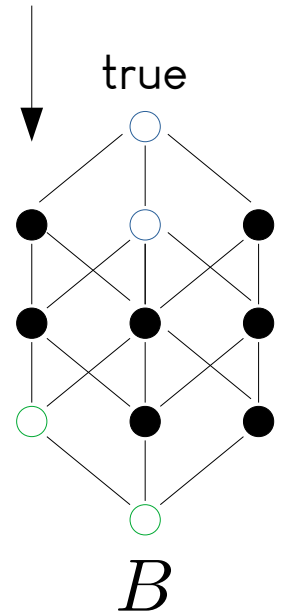
Query  
Elise: 8h - 12h  
Paul: 10h - 23h



no

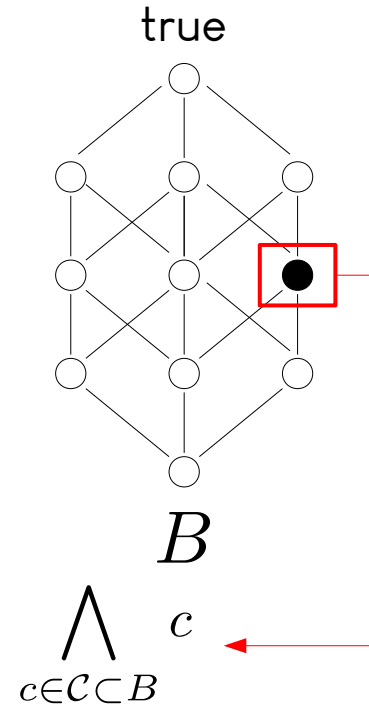
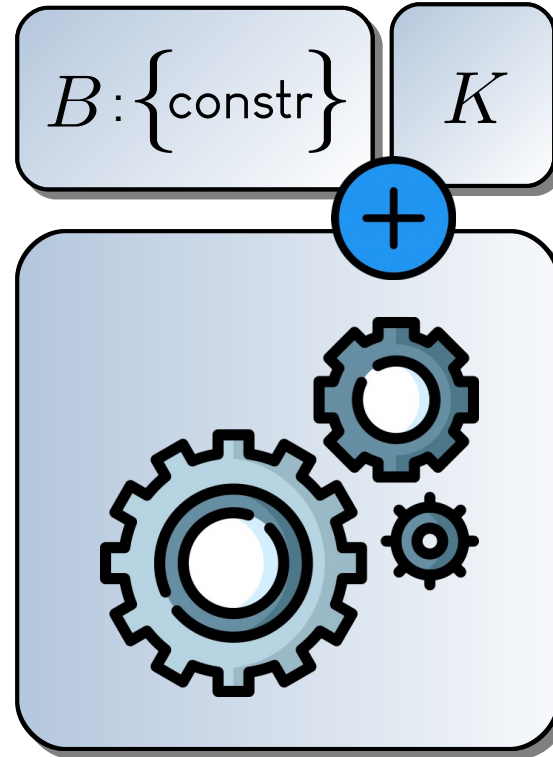
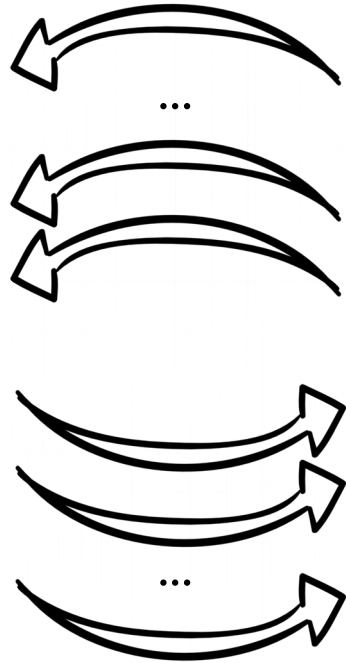


no: Top-down

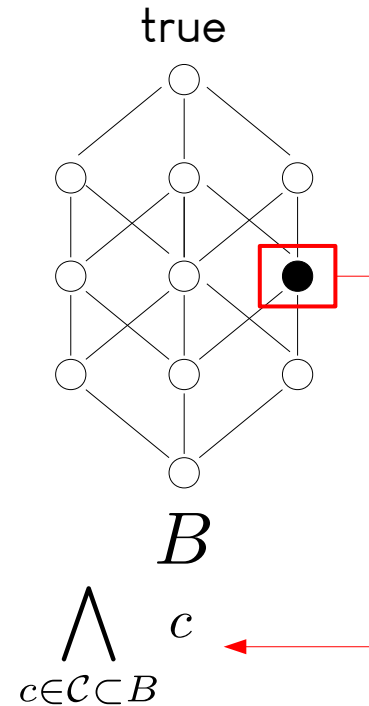
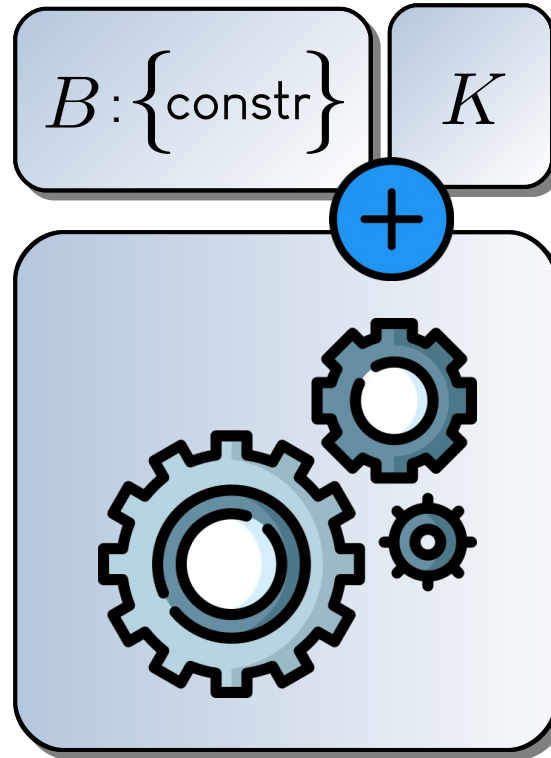
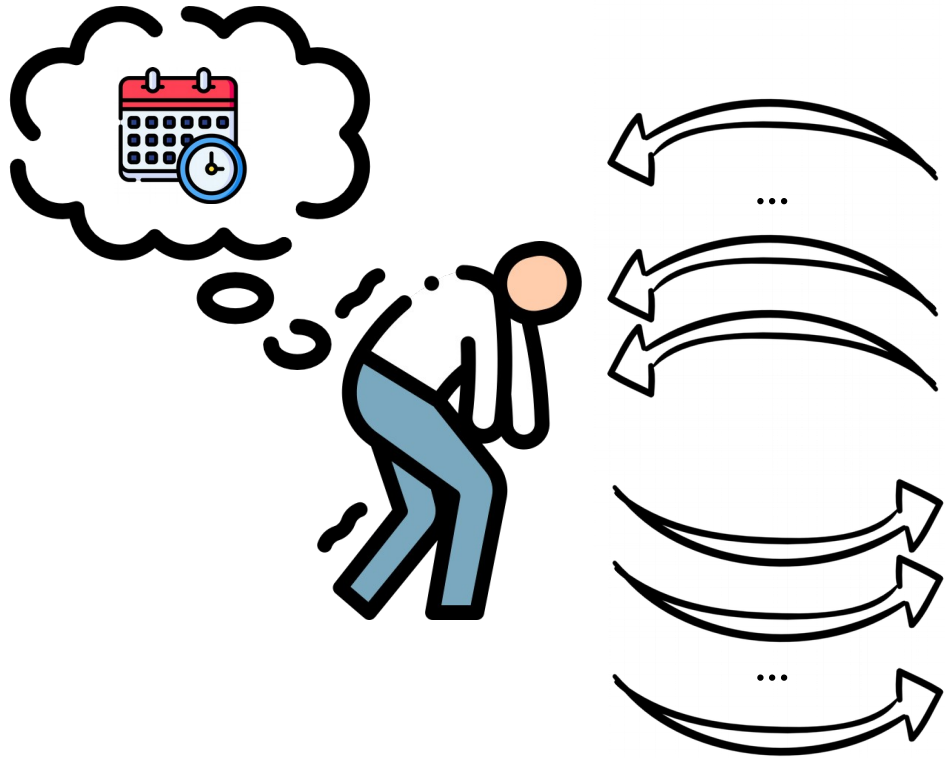




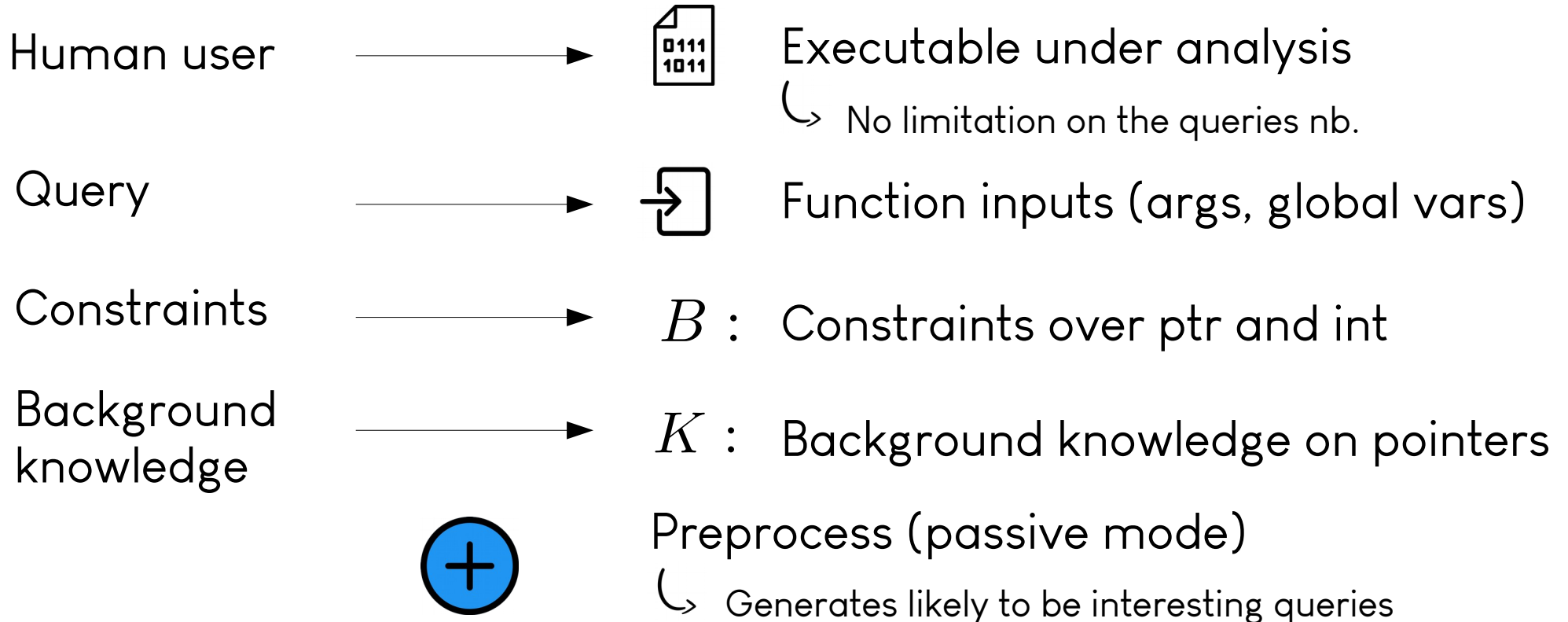
# Active Constraint Acquisition



# Careful: too many queries



# Adapting Constraint Acquisition



# Which constraints ?

Constraints  $\longrightarrow$   $B$  : Constraints over ptr and int

Constraints for memory-related precondition.:

$P$   $:=$   $C \Rightarrow A \mid A \mid \neg A$

$C$   $:=$   $C \wedge C \mid A \mid \neg A$

$A$   $:=$   $valid(p_j) \mid alias(p_j, p_l) \mid deref(p_j, g)$   
 $\mid i_j = 0 \mid i_j < 0 \mid i_j \leq 0 \mid i_j = i_l \mid i_j < i_l \mid i_j \leq i_l$

Method not limited to  
memory-related precondition.



**From language to bias B:** max. size of horn clauses depending on the function prototype – especially number of integer inputs

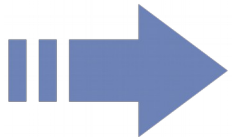
# Preprocess

## Positive ( $e^+$ ) vs Negative Queries ( $e^-$ )

↳ **All constraints** incoherent with  $e^+$  are not in the solution



↳ **At least one constraint** incoherent with  $e^-$  is in the solution



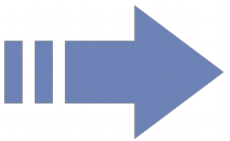
Generate positive queries first to remove a lot of constraints

↳ How to generate positive queries ?

# Preprocess



- Goal of developers : software should work
  - Usually they handle well usual cases
- Generate queries where code likely behave correctly



Generate first queries with  $\leq 1$  one non valid, aliasing or deref pointers

# PreCA

NEW

Call the preprocess

```
while true do
```

```
  Generate an informative query
```

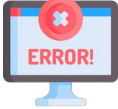


```
  if no-query then «we converged»
```

```
  Submit query to the oracle(F, Q)
```

```
  if answer is yes then  
  | Bottom-up-inference()
```

```
  else  
  | Top-down-inference()
```

## How Oracle answers queries ?

- ↳ Run function F under query
- ↳ If  $ret \neq Q$  or  .....▶ no
- ↳ If  .....  .....▶ ukn
- ↳ Otherwise .....▶ yes



# Back To Our Example



```
find_first_of(int* a, int m, int* b, int n)
```

Description: returns the index of the first element in “a” present in “b”

Postcondition:  $Q = true$



Variables : a, m, b, n



Heuristics : max. Horn clause size = 3

# Back To Our Example



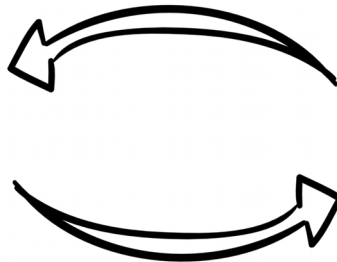
Fun: `find_first_of`  
In : `int* a, int m,`  
`int* b, int n`

Description: returns the index of the first element in “a” present in “b”

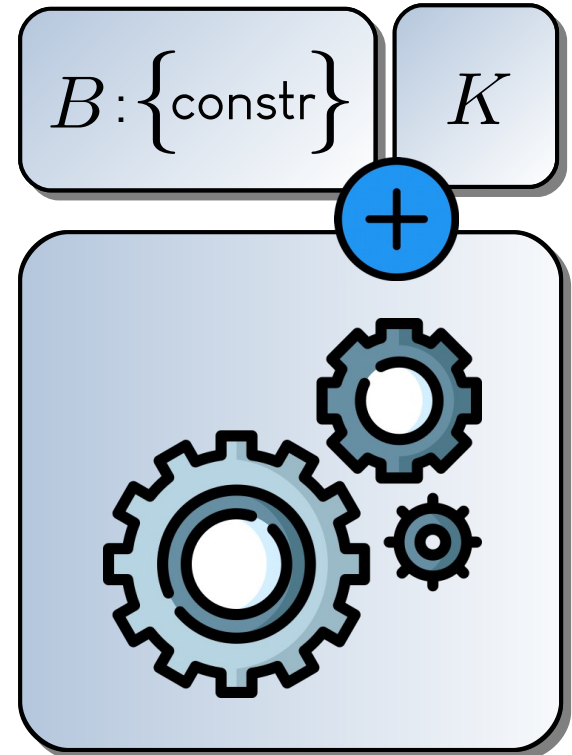
Postcondition:  $Q = true$

Query

`a = NULL, m = 1`  
`b = NULL, n = 3`



no



# Back To Our Example



Fun: `find_first_of`  
In : `int* a, int m,`  
`int* b, int n`

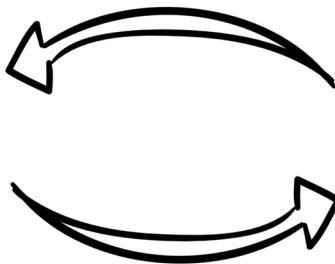
Description: returns the index of the first element in “a” present in “b”

Postcondition:  $Q = true$

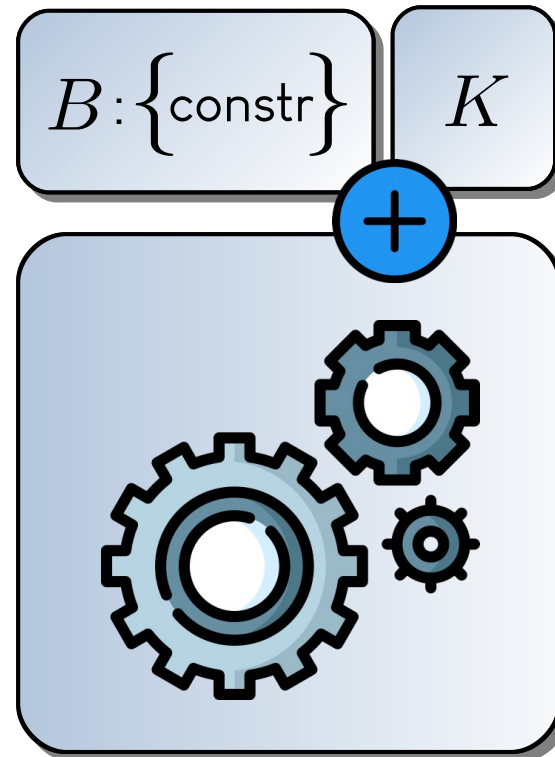
Query

$a = NULL, m = 0$

$b = NULL, n = 0$



yes



# Back To Our Example



```
Fun: find_first_of  
In : int* a, int m,  
     int* b, int n
```

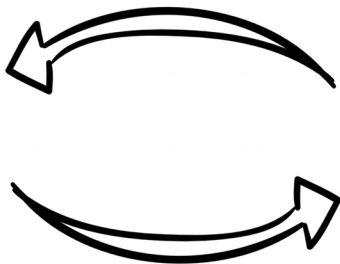
Description: returns the index of the first element in “a” present in “b”

Postcondition:  $Q = true$

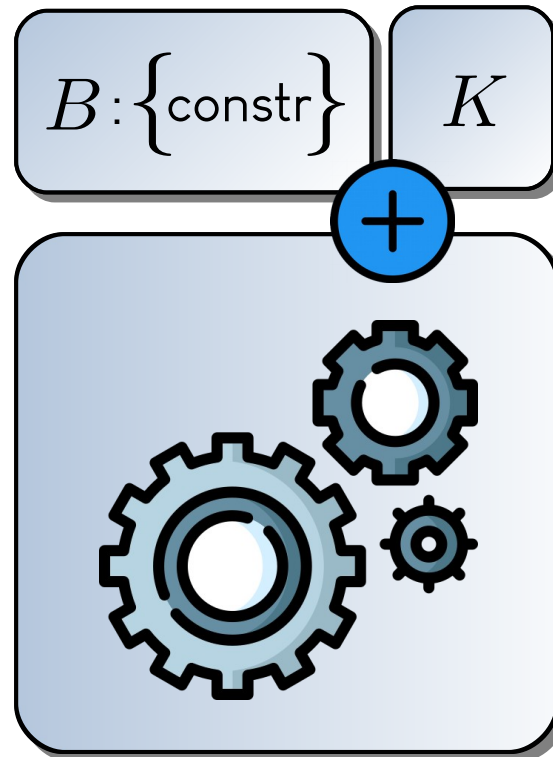
Query

$a = @1, m = 3$

$b = @2, n = 3$



yes



# Back To Our Example



Fun: `find_first_of`  
In : `int* a, int m,`  
`int* b, int n`

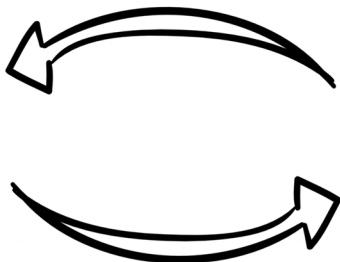
Description: returns the index of the first element in “a” present in “b”

Postcondition:  $Q = true$

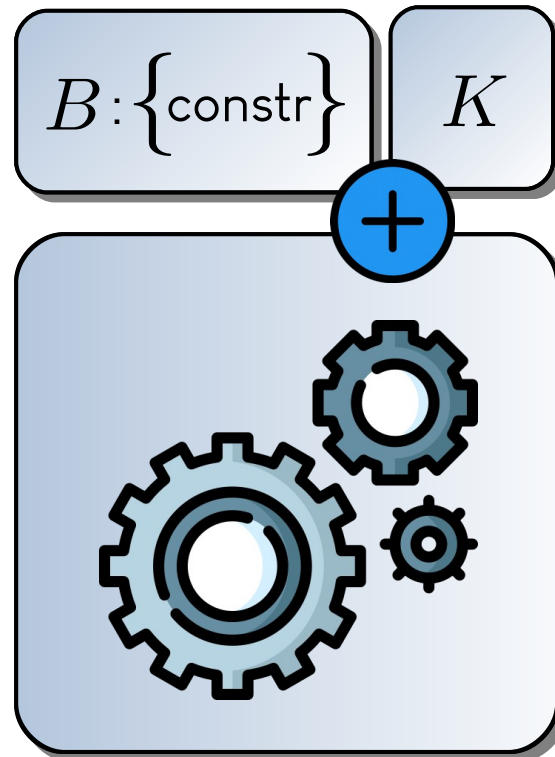
Query

$a = @1, m = 3$

$b = @1, n = 3$



yes



# Back To Our Example

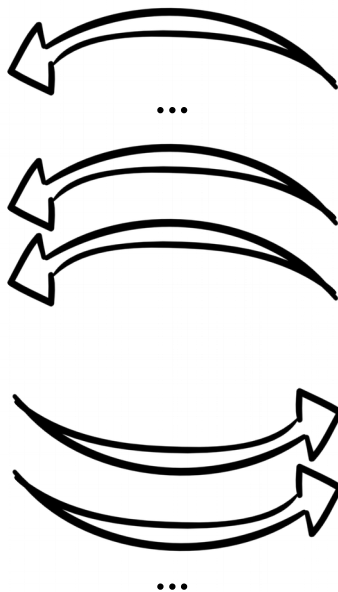


```
Fun: find_first_of  
In : int* a, int m,  
     int* b, int n
```

Description: returns the index of the first element in “a” present in “b”

Postcondition:  $Q = true$

45 queries



$B: \{ \text{constr} \}$

$K$

Result



$[m > 0 \Rightarrow \text{valid}(a)]$

$\wedge$

$[m > 0 \wedge n > 0 \Rightarrow \text{valid}(b)]$

# Theoretical Analysis

PreCA guarantees

- ↳ If B is expressive enough .....  or Precond.
- ↳  If oracle never answers “unk” ..... The most general precondition

These are good theoretical guarantees

- ↳ SOTA executions based methods, from programming language community, have no clear guarantees



# Evaluation

**Dataset:** 94 learning tasks • compiled C functions (string.h, arrays, arithmetic ...)

**Evaluation:** \_\_\_\_\_

1 hour

PreCA

92%

41%



Daikon, PIE, Gehr et al

At most 52%

At most 23%



P-Gen

74%

34%

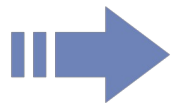


PreCA better in 5s than concurrent tools in 1 hour

# Conclusion

## AI contributions

- ↳ 1<sup>st</sup> adaptation of CA for prog. analysis
  - new use case for CA
  - no user (no limit for queries nb)
- ↳ Translate core concepts :
  - Set of constraints
  - Background knowledge
- ↳ Extend CA (ukn, preprocess)



**Opens new research directions for CA**

## Prog. analysis contribs

New efficient precond. inference tool



Good guarantees



Outperforms concurrent tools



Does not need the source code

Thank you for your  
attention



@grmenguy



<https://gregoiremenguy.github.io/>